

Generating Random Variates I

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Course overview

1. Introduction

7. NS-2: Fixed networks

2. Building block: RNG

8. NS-2: Wireless networks

**3. Building block:
Generating random variates I
and modeling examples**

9. Output analysis: single system

**4. Building block:
Generating random variates II
and modeling examples**

**10. Output analysis: comparing
different configuration**

**5. Algorithmics:
Management of events**

11. Omnet++ / OPNET

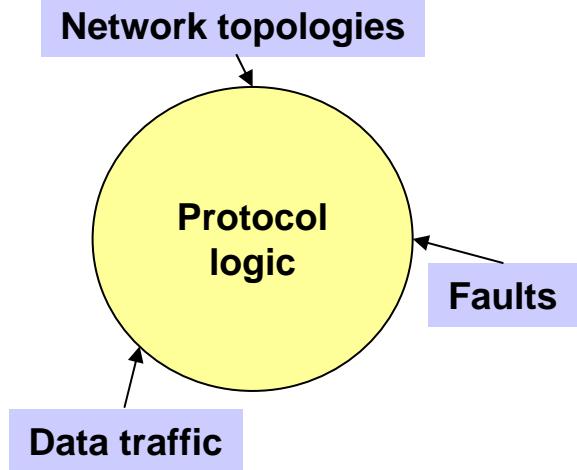
6. NS-2: Introduction

12. Simulation lifecycle, summary

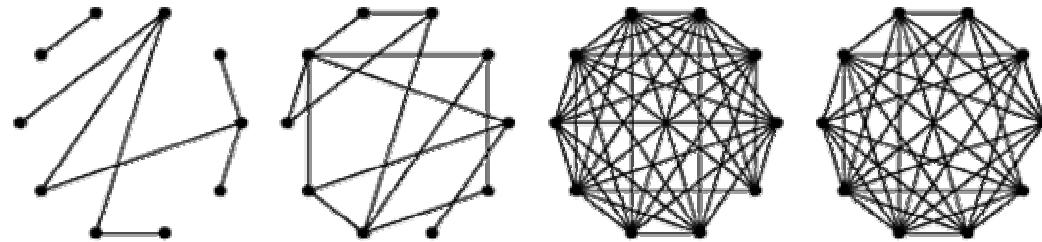
Structure of this lecture

- » **Part I: Modeling examples w.r.t. $U(0,1)$**
 - Random graphs (topologies)
 - Constant bit rate data traffic with jitter
 - Randomness as protocol element
 - Example: multiaccess communication
- » **Part II: Generation of other random variates for modeling network elements**
 - Inverse transform method
 - Example: exponential distribution
 - Inverse transform method for discrete random variates
 - Generalized inverse transform
- » **Part III: Modeling examples w.r.t. exponential distribution**
 - Random direction mobility
 - On/off sources for generating bursty traffic

Random effects:
'how to' and 'what'



I Random graphs I



Eric W. Weisstein. "Random Graph." From [MathWorld](http://mathworld.wolfram.com/RandomGraph.html)--A Wolfram Web Resource.
<http://mathworld.wolfram.com/RandomGraph.html>

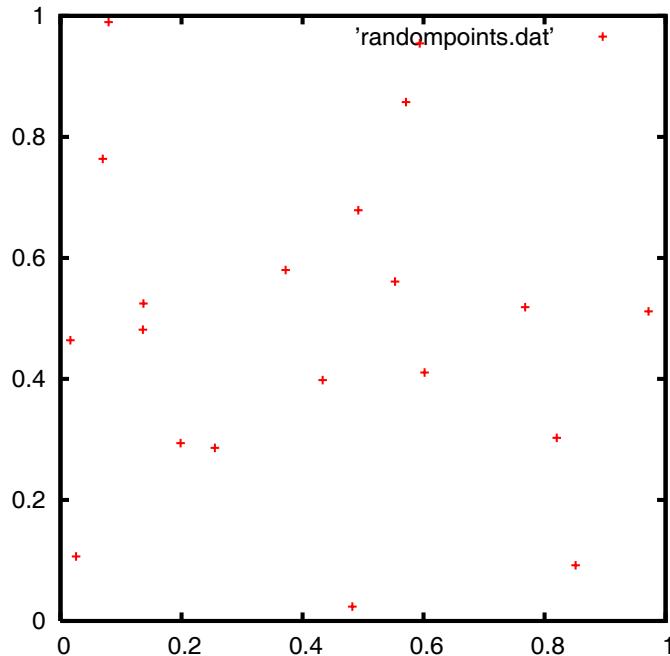
For each pair of vertices (u,v) :

- » **There is an edge between u and v with probability $a \in [0,1]$.**
- » **Required: sampling from Bernoulli random variable; mass function:**

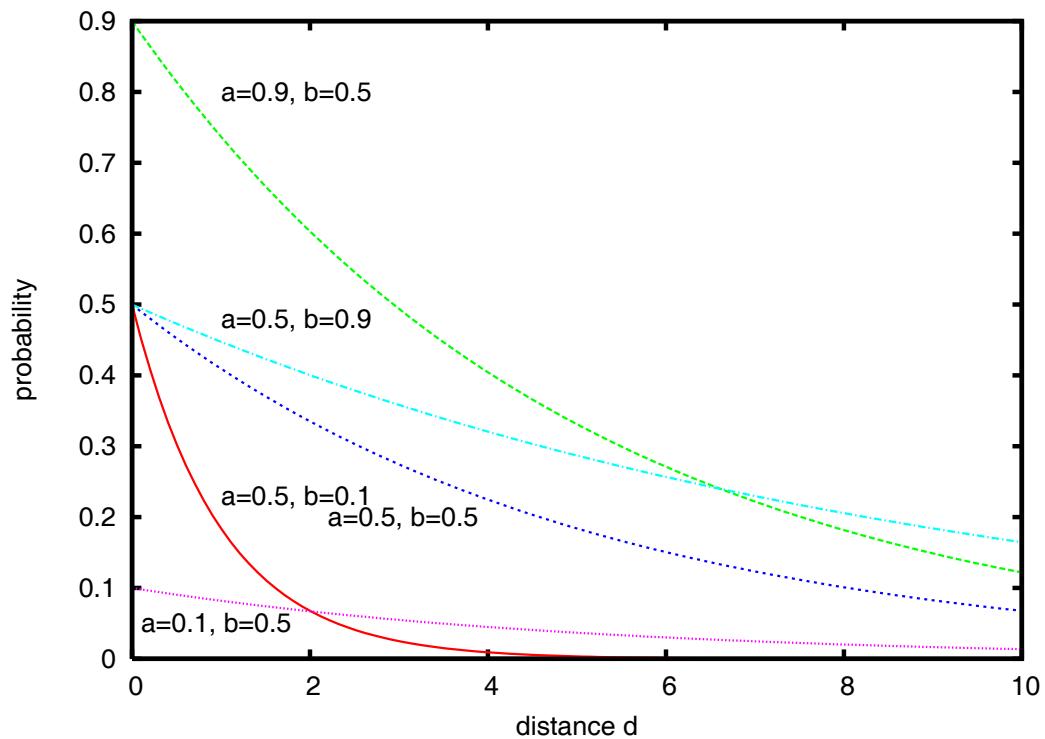
$$p(x) = \begin{cases} 1 - a & \text{if } x = 0 \\ a & \text{if } x = 1 \\ 0 & \text{otherwise} \end{cases}$$

I Random graphs II

- » **Select vertices by sampling from a 2-dimensional uniform distribution**
 - Since dimensions should be independent, one can simply sample x and y coordinates from a 1-dimensional uniform distribution



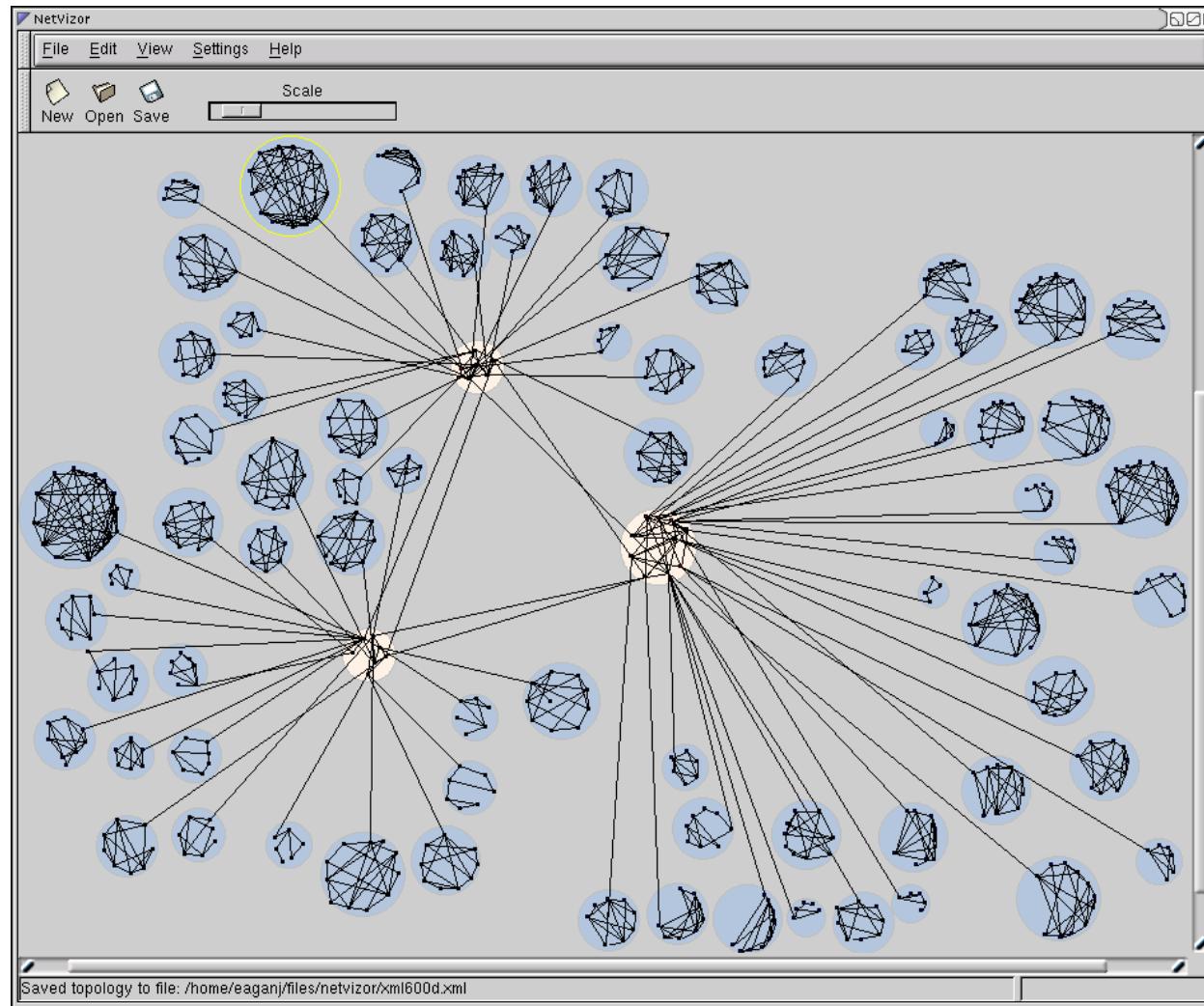
I Random graphs III: Waxman model



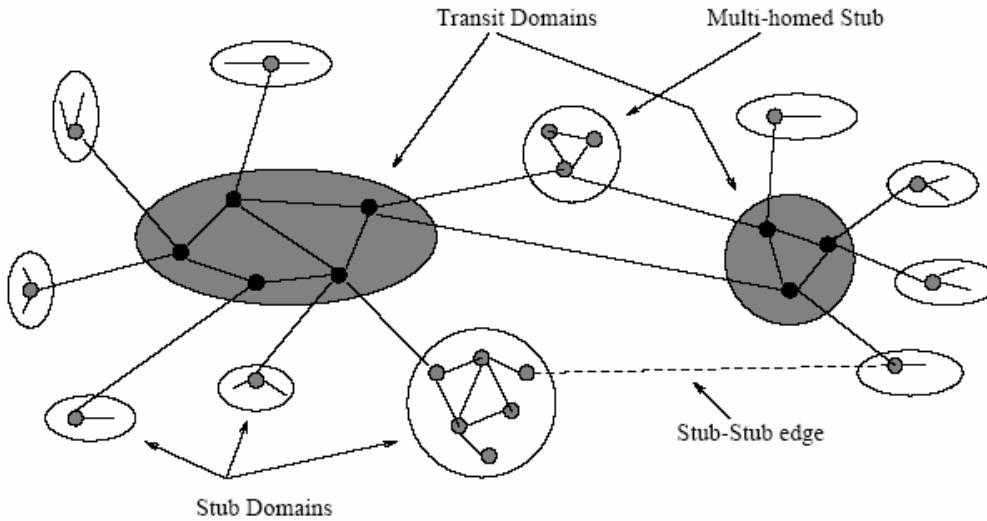
$$P(u, v) = a \exp\left(\frac{-d}{bL}\right)$$

- » Caution: this is not an exponential distribution. It is a Bernoulli one (edge yes or no) where the probability of ‘edge’ depends on the Euclidean distance d of the new nodes u and v , and on the system areas diameter L .
- » Parameter a adjusts the decay w.r.t. distance, b the general level of ‘success’.

I Hierarchical random graphs



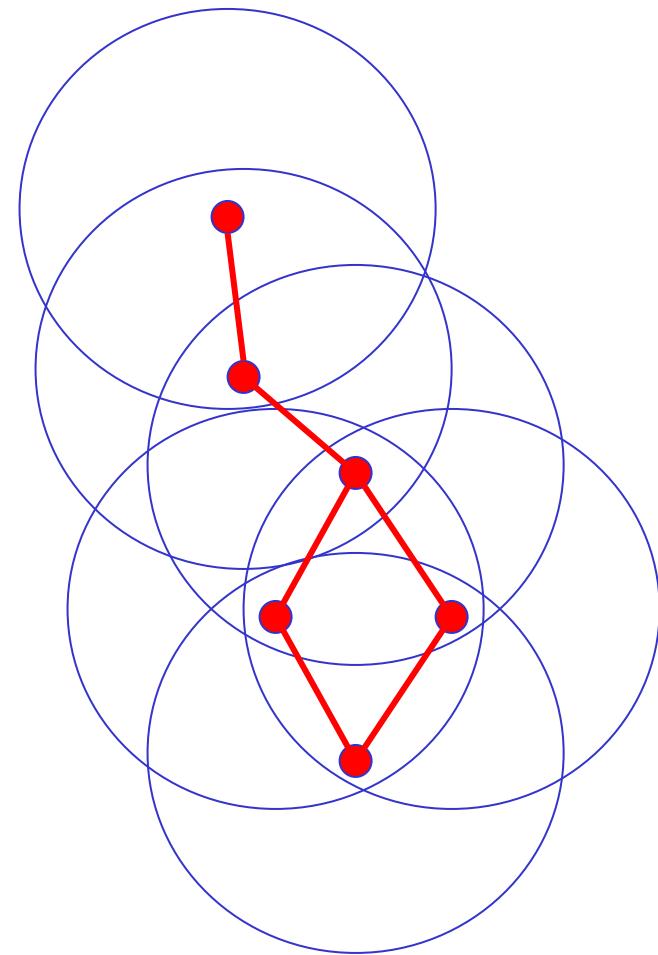
I Hierarchical random graphs



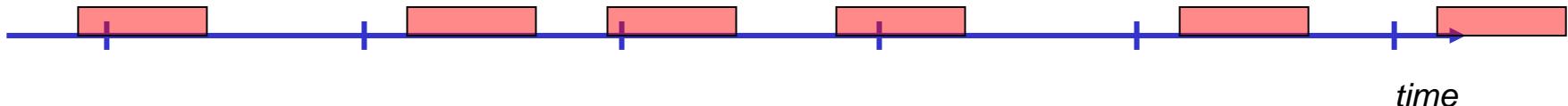
- » **Construct connected random graph: each node represents an entire transit domain**
- » **Each node is replaced by a connected random graph representing a transit network**
- » **For each node in transit networks: construct connected random graph representing the associated stub networks**
- » **Add some more edges between transit-and-stub and stub-and-stub networks.**

I Unit disc graph (for wireless scenarios)

- » For wireless scenarios – idealized ‘transmission ranges’
- » There is an edge between nodes u and v if distance between nodes is smaller than a given value t .



I Constant bit-rate traffic with jitter

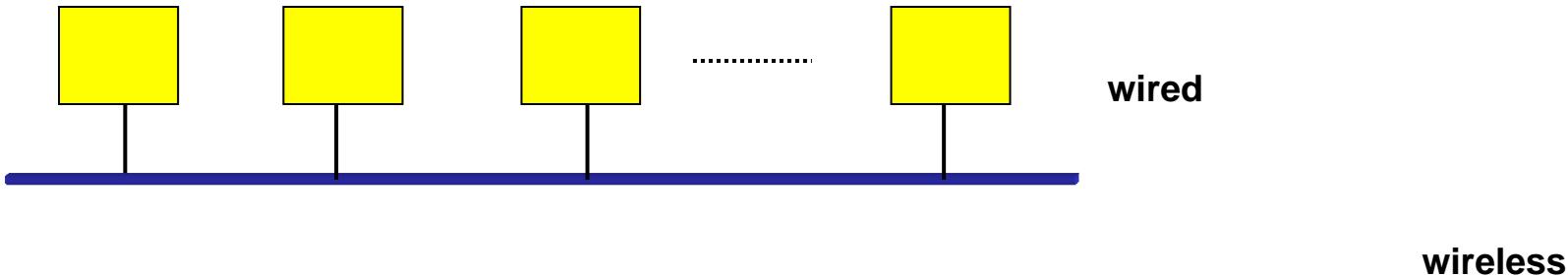


- » Constant bit rate traffic is parameterized via
 - Rate: the sending rate or
 - Interval: interval between packets and
 - PacketSize: the constant size of the packets generated
- » Introducing jitter: random flag indicating whether or not to introduce random “noise” in the scheduled departure times (default is off)
- » Example:
 - packetSize 48
 - rate 64Kb
 - random 1
- » Example (NS-2, /tools/cbr_traffic.cc):

```
double t = interval_; if (random_)  
    t += interval_ * Random::uniform(-0.5, 0.5);
```

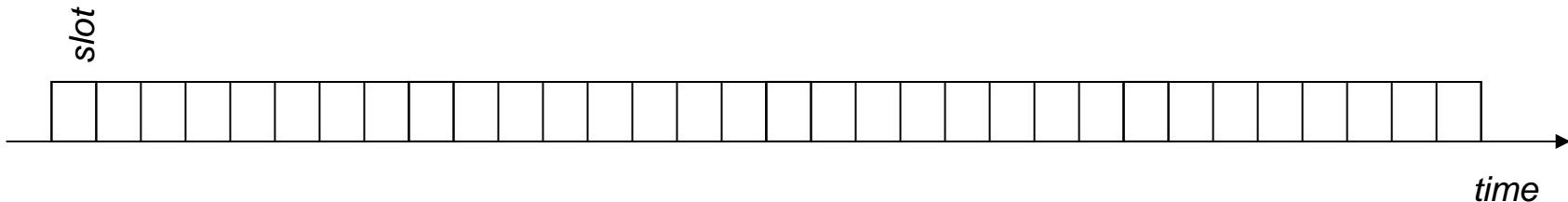
I Multiaccess communication: problem statement

- » Randomness as a mechanism for self-organization

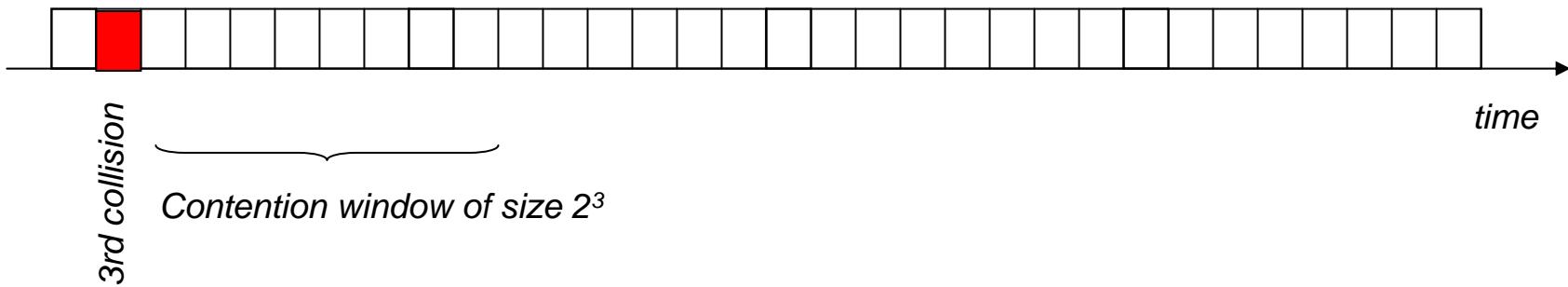


- » Assume: no centralized control
- » How to organize access to shared medium?
- » When two or more stations send a packet at the 'same' time, a collision occurs

I Multiaccess communication: exponential backoff



- » **Binary exponential backoff:**
- » If a packet has been transmitted unsuccessfully i times, the transmission is rescheduled for a randomly chosen slot in the ‘contention window’ of size 2^i slots.
- » Self-adjustment to number of competing stations.
- » Example:



I Multiaccess communication: 10Mbps Ethernet

Collision on attempt no.	Contention window	Backoff times
1	0 ... 1	0 ... 51.2 μ s
2	0 ... 3	0 ... 153.6 μ s
3	0 ... 7	0 ... 358.4 μ s
4	0 ... 15	0 ... 768 μ s
5	0 ... 31	0 ... 1.59 ms
6	0 ... 63	0 ... 3.23 ms
7	0 ... 127	0 ... 6.5 ms
8	0 ... 255	0 ... 13.1 ms
9	0 ... 511	0 ... 26.2 ms
10-15	0 ... 1023	0 ... 52.4 ms
16	N/A	Discard frame

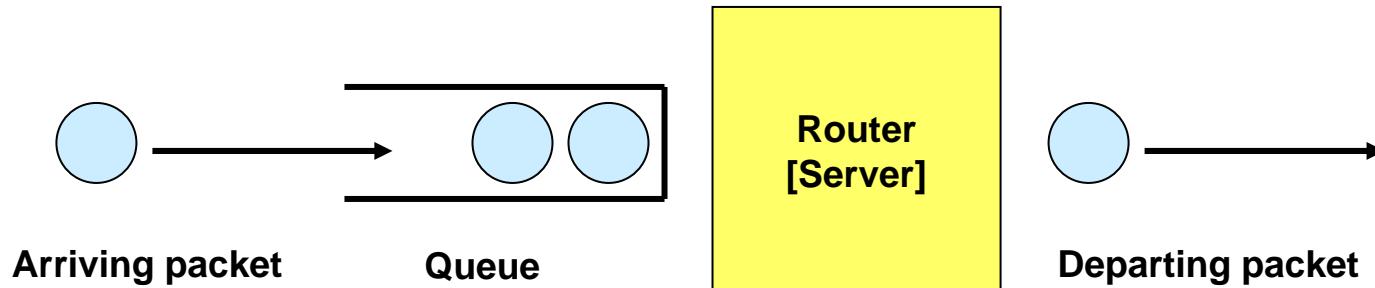
I Multiaccess communication: NS-2 example

```
rtime = (Random::random() % cw) * mac->phymib_.getSlotTime();
```

- » From ns-2.27/mac/mac-timers.cc
- » random() returns random integer in [0, MAXINT]
- » Efficient?
- » Other approach starting with U(0,1):

$$\lfloor (cw + 1)U \rfloor$$

II Generating other random variates: introduction



How to generate exponentially distributed inter-arrival times?

- » Algorithms to produce observations (“variates”) from some desired input distribution (exponential, gamma, etc.)
- » Formal algorithm - depends on desired distribution.
- » But *all* algorithms have the same general form:
 - Generate one or more IID $U(0, 1)$ random numbers
 - Transformation (depends on desired distribution)
 - Return $X \sim$ desired distribution
- » Note critical importance of a good random-number generator

II Criteria

May be several algorithms for a desired input distribution form; want:

- » **Exact:** X has *exactly* (not approximately) the desired distribution
 - Example of approximate algorithm:
 - Treat $Z = 6 - U_1 + U_2 + \dots + U_{12}$ as normal distribution $N(0,1)$
 - Mean, variance correct; rely on Central Limit Theorem for *approximate* normality
 - Range clearly incorrect
- » **Efficient:**
 - Low storage
 - Fast (marginal, setup)
 - Efficient regardless of parameter values (*robust*)
- » **Simple:** Understand, implement (often tradeoff against efficiency)

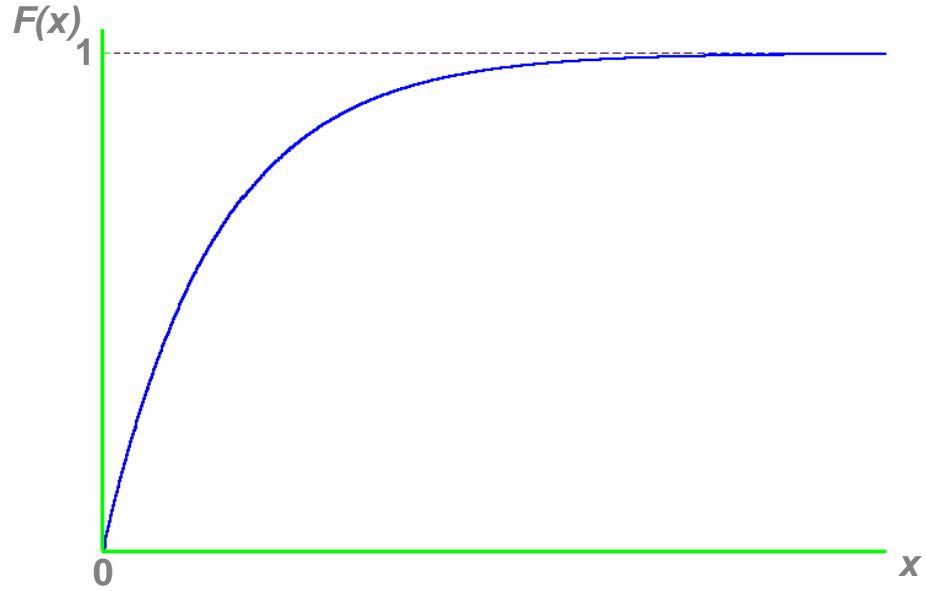
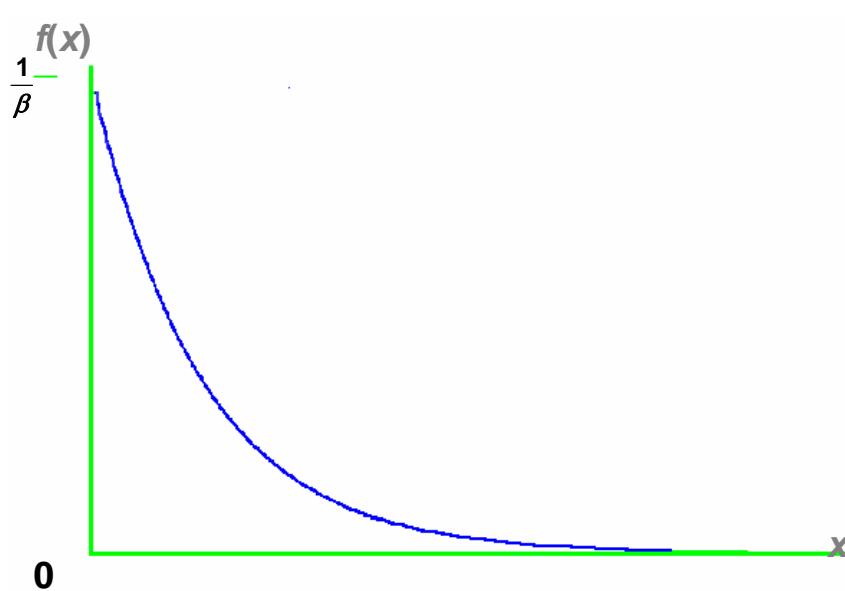
II Recap: distribution function

Definition:

the (cumulative) distribution function $F(x)$ for a random variable

$X: \Omega \rightarrow \mathbb{R}$ is defined as

$$F(x) = P(X \leq x) \quad \text{for} \quad -\infty < x < \infty$$



II Recap: strictly increasing function

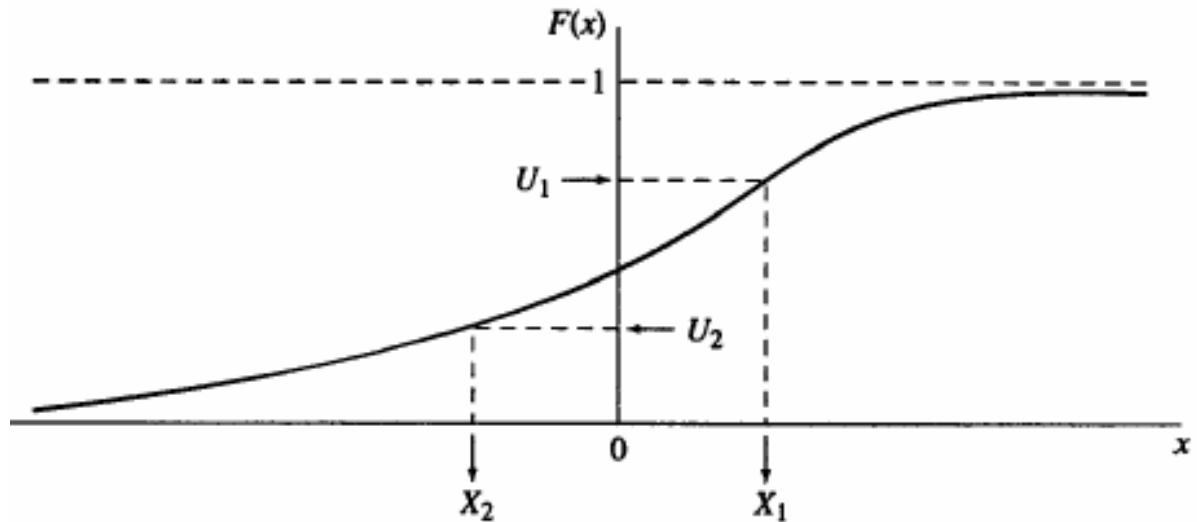
- » A distribution function F is *strictly increasing* when

$$x_1 < x_2 \rightarrow F(x_1) < F(x_2)$$

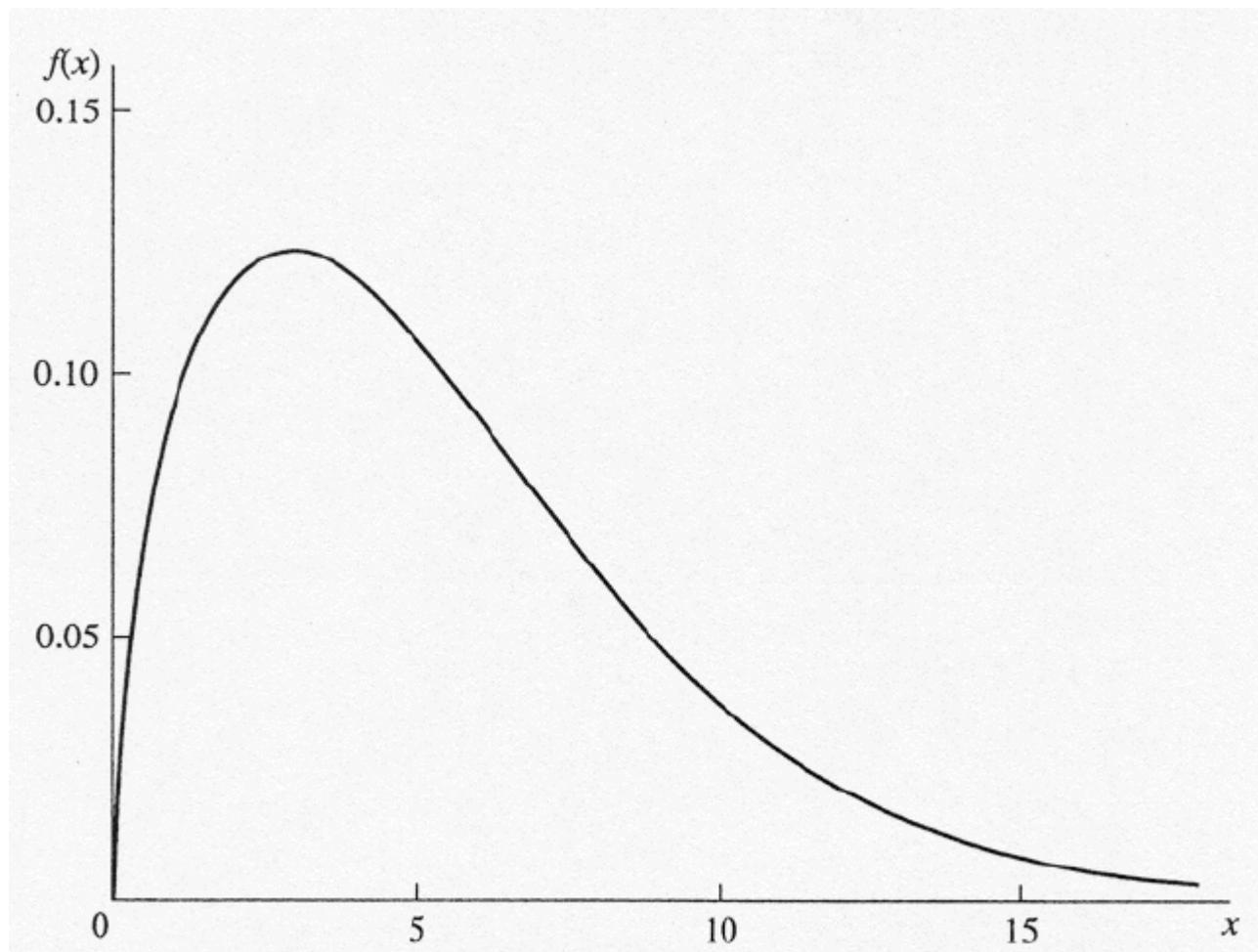
- » For a distribution function F that is continuous and strictly increasing when $0 < F(x) < 1$, the inverse function F^{-1} exists on $(0,1)$.

II Inverse transform method

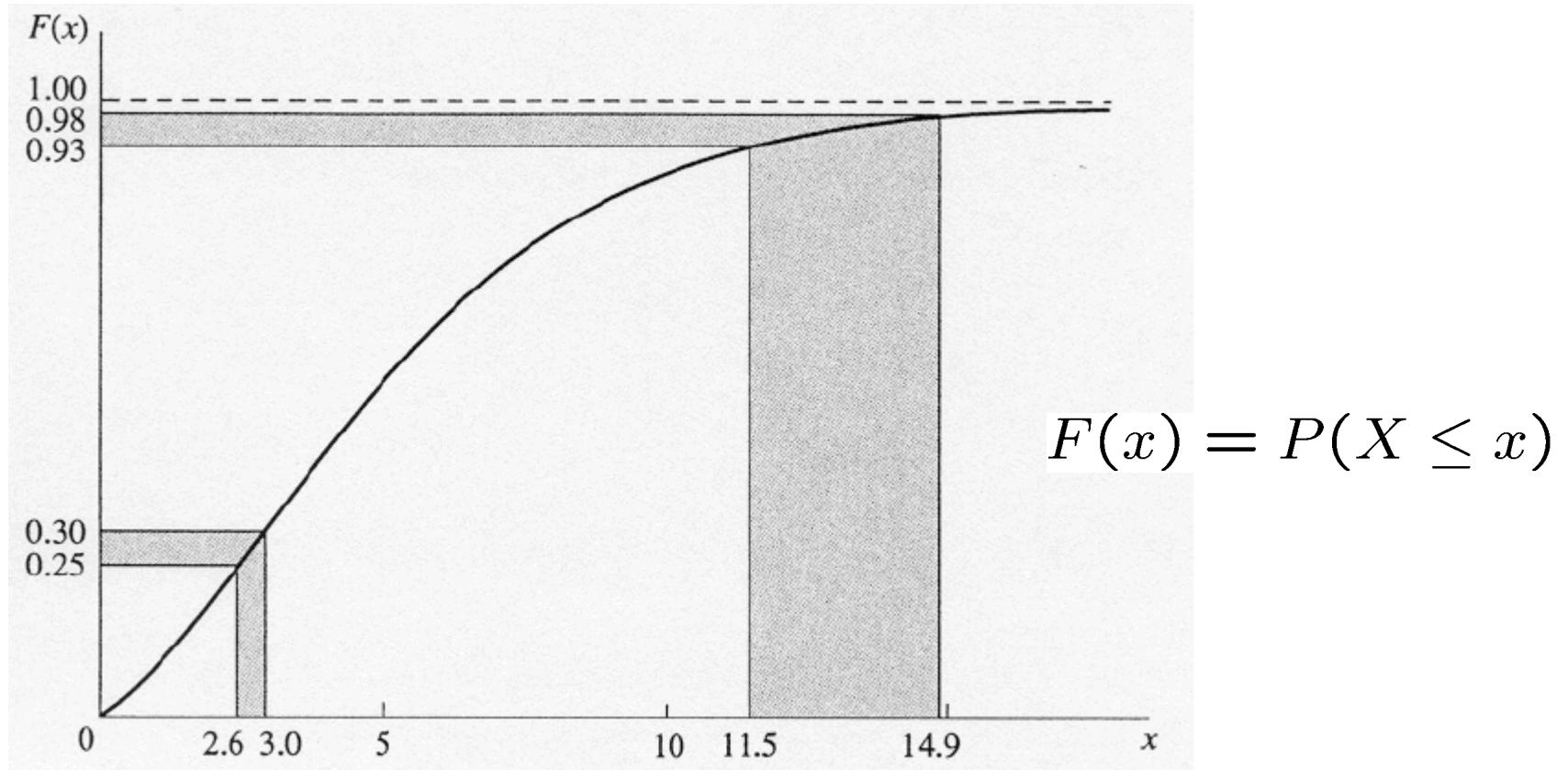
- » Goal: generate a random variate X that is continuous and has distribution function F that is continuous and strictly increasing when $0 < F(x) < 1$.
- » Strategy: Inverse transform algorithm
 1. Generate $U \sim U(0,1)$
 2. Return $X = F^{-1}(U)$



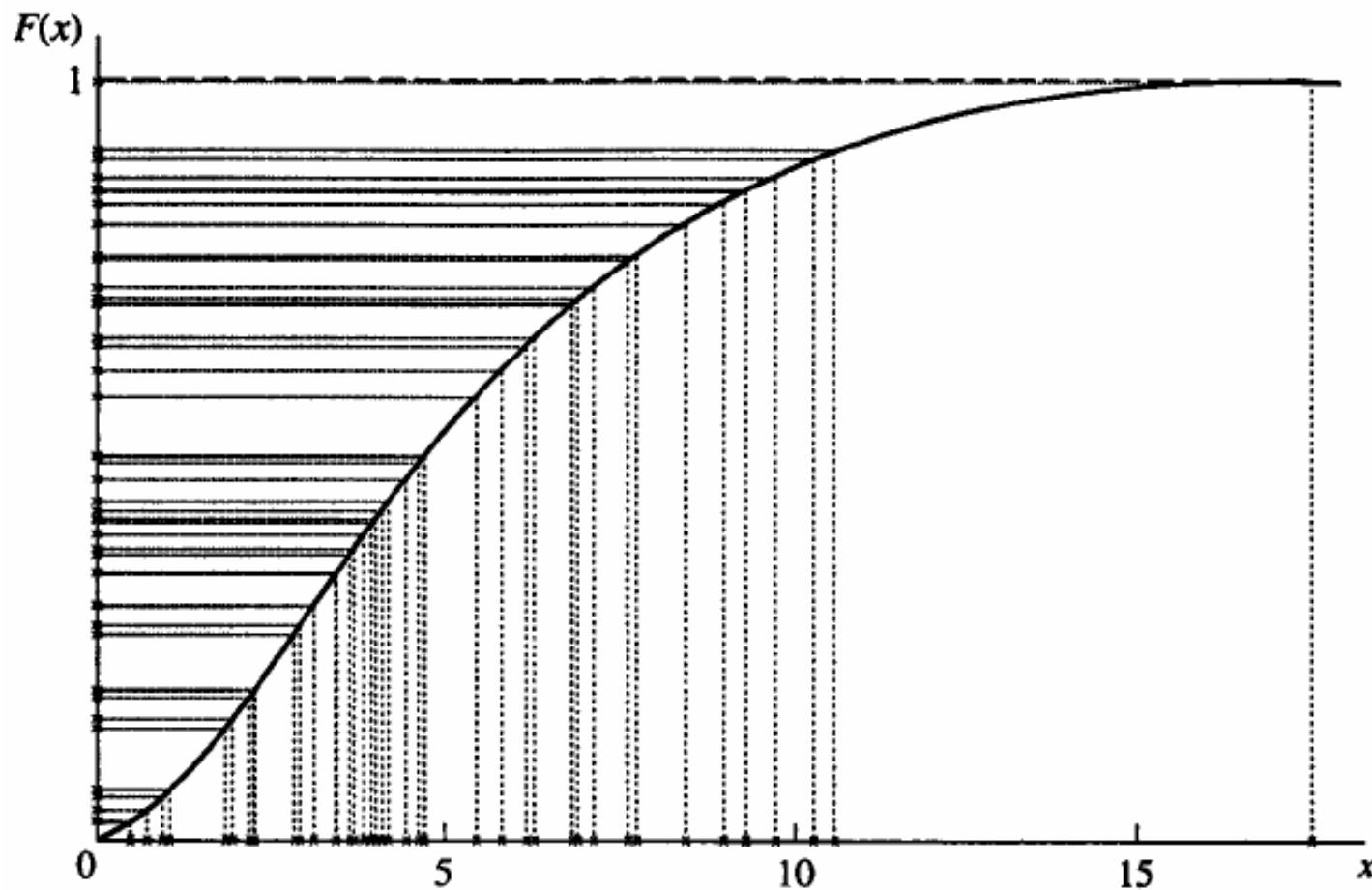
II Intuitive explanation I: density



II Intuitive explanation II: distribution function



II Intuitive explanation III: algorithm at work (50 samples)



II Proof

$$P(X \leq x) = P(F^{-1}(U) \leq x)$$

» **Definition**

$$P(F^{-1}(U) \leq x) = P(U \leq F(x))$$

» **Properties of F**

$$P(U \leq F(x)) = F(x)$$

» **Uniform distribution,
definition of distribution
function**

$$P(X \leq x) = F(x)$$

II Example: exponential distribution

» Density

$$f(x) = \begin{cases} \frac{1}{\beta} e^{-x/\beta} & \text{if } x \geq 0 \\ 0 & \text{otherwise} \end{cases}$$

» Distribution function

$$F(x) = \begin{cases} 1 - e^{-x/\beta} & \text{if } x \geq 0 \\ 0 & \text{otherwise} \end{cases}$$

» Inverse transform

$$F^{-1}(u) = -\beta \ln(1 - u)$$



uniformly distributed

II Exponential distribution: C code

From M/M/1 example:

```
float expon(float mean) /* Exponential variate
    generation function. */

{/* Return an exponential random variate with mean
 "mean". */

    return -mean * log(lcgrand(1));}
```

From NS-2 (tools/rng.h):

```
inline double exponential ()

{return (-log(uniform()));}

inline double exponential (double r)

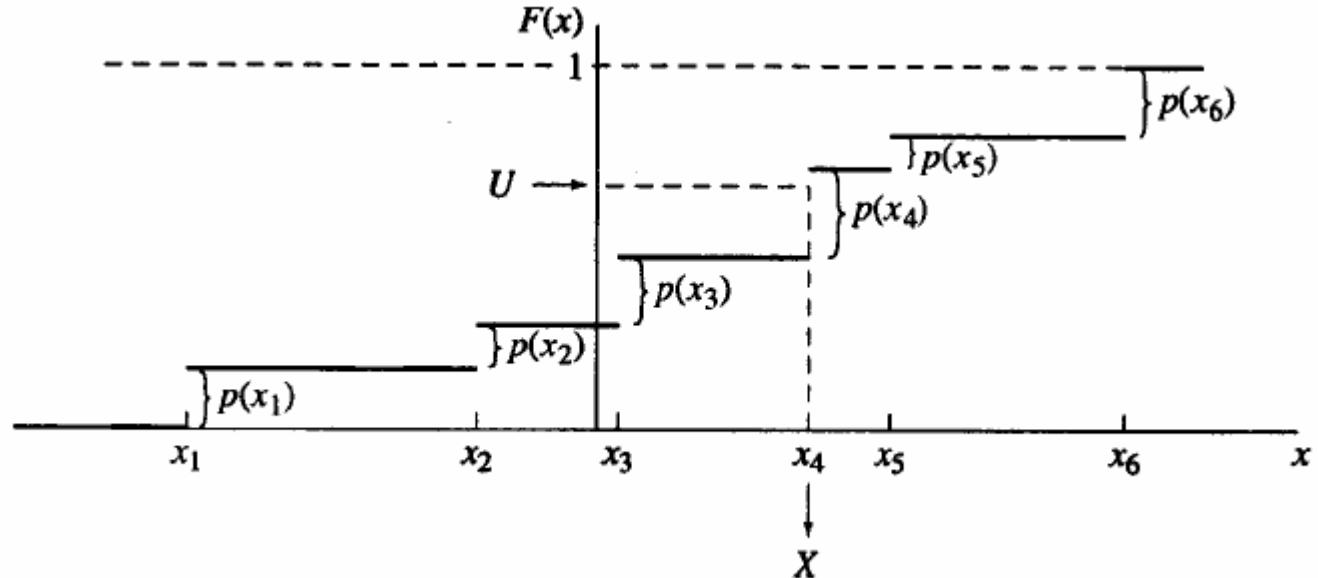
{return (r * exponential());}
```

II Inverse transform method for discrete distributions

- » Assume: random variable X can take only values x_1, x_2, \dots with $x_1 < x_2 < \dots$
- » Probability mass function: $p(x_i) = P(X = x_i)$
- » Distribution function: $F(x) = P(X \leq x) = \sum_{x_i \leq x} p(x_i)$

Algorithm:

1. Generate $U \sim U(0,1)$
2. Determine the smallest positive integer I such that $U \leq F(x_I)$, and return X



II Discrete case: efficiency considerations

Step 2 involves a “search” of some kind; several computational options:

- Direct left-to-right search—if $p(x_i)$ ’s fairly constant
- If $p(x_i)$ ’s vary a lot, first sort into decreasing order, look at biggest one first, ..., smallest one last—more likely to terminate quickly
- Example:
 - Probabilities $\{p_1=1/8, p_2=1/8, p_3=1/4, p_4=1/2\}$.
 - A: order 1, 2, 3, 4. Expected number of comparisons: $11/8$
 - B: order 4, 3, 2, 1. Expected number: $9/8$
- Exploit special properties of the form of the $p(x_i)$ ’s, if possible
 - To facilitate computation of $F(x)$



II Generalizations and limits

» Generalized Inverse-Transform Method

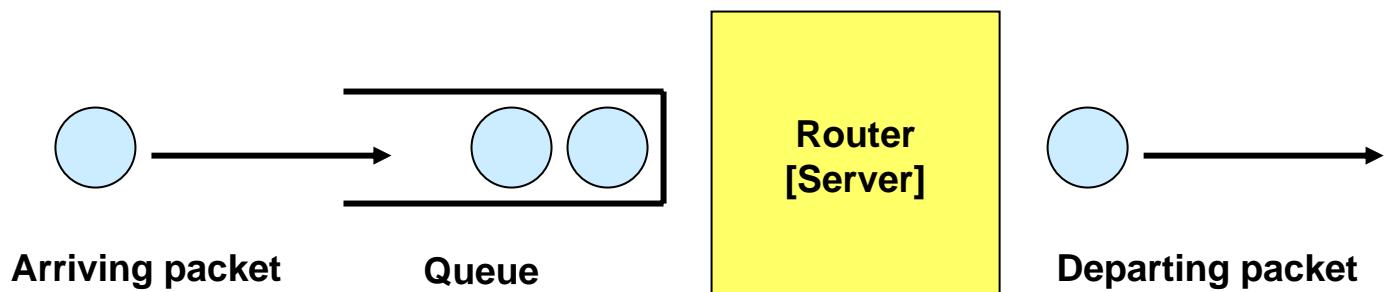
- Valid for *any* CDF $F(x)$: return $X = \min\{x: F(x) \geq U\}$, where $U \sim U(0,1)$
 - Continuous, possibly with flat spots (i.e., not strictly increasing)
 - Discrete
 - Mixed continuous-discrete

» Problems with Inverse-Transform Approach

- Must invert CDF, which may be difficult (numerical methods)
- May not be the fastest or simplest approach for a given distribution

III Modeling examples

- » ... using exponential variates
- » Recap: Inter-arrival times



- » Examples for topologies
 - Mobile case: random direction mobility
- » Examples for synthetic data traffic
 - On/off sources for generating bursty traffic

III Random direction mobility

- » Last lecture: random waypoint mobility
- » Various flavors around
- » Node chooses random direction sampled from uniform distribution over $[0, 2\pi)$
- » Node chooses speed either from uniform or normal distributions
- » Node changes direction/speed after a time interval sampled from an exponential distribution
- » At system area's boundary: wrap-up, reflection ...

III On/off sources for modeling bursty traffic

- » NS-2: EXPOO_Traffic—generates traffic according to an Exponential On/Off distribution. Packets are sent at a fixed rate during on periods, and no packets are sent during off periods. Both on and off periods are taken from an exponential. Packets are constant size.
- » **packetSize_** the constant size of the packets generated
- » **burst_time_** the average “on” time for the generator
- » **idle_time_** the average “off” time for the generator
- » **rate_** the sending rate during “on” times
- » Example:
 - **packetSize_** 210
 - **burst_time_** 500ms
 - **idle_time_** 500ms
 - **rate_** 100k

III VoIP modeling

- » ITU-T recommendations for VoIP modeling:
- » On periods: exponential with mean 1.004s
- » Off periods: exponential with mean 1.587s
- » Sending rate during on periods: 88kbps

Summary

- » Examples of how to model network topologies and data traffic
- » Discussion on how to generate arbitrary random variates
- » Inverse transform method
 - Particularly good for exponential variates

So far:

- Generation of uniform, discrete and exponential variates
- Random topologies (static, wireless, mobile networks),
Synthetic data traffic (CBR, exponentially distributed
inter-arrival times, on/off-sources)

References

- » Exponential backoff
 - Charles E. Spurgeon, *Ethernet – The Definite Guide*, O'Reilly, 2000; pp. 53-60
 - Dimitri Bertsekas, Robert Gallager, *Data Networks*, 2nd ed., Prentice Hall, 1992, pp. 286
- » Inverse transform method
 - Averill M. Law, W. David Kelton, *Simulation Modeling and Analysis*, 3rd ed., McGrawHill, 2000
- » Waxman model
 - Bernard M Waxman Routing of multipoint connections IEEE Journal on Selected Areas in Communications, 6 (9): 1617-1622, 1988.
- » Hierarchical random graphs
 - Georgia Tech Internet Topology Map (GT-ITM), *models.ps*